#### Johannes Fischer Advanced Data Structures

MSc-Vorlesung Wintersemester 2012/13 KIT

#### Preliminaries

- 5 ECTS
- lectures in German, slides etc. in English
- prequisites:
  - Algorithmen II
  - interest in discrete, combinatorial problems
- ~14 lectures (NOT 27.12.12/03.01.13)
- oral exam (20-25 mins)

#### Preliminaries

- course homepage: http://algo2.iti.kit.edu/2056.php
  - slides & script
  - additional course information
- Johannes.Fischer@kit.edu (room 207)
- office hours: Thursday 14-15

# BADS'I 3

- write a **paper** on a data structure **not** covered in the lecture
  - list of topics on course website
  - can be strengthened by experiments
- use LaTeX  $\Rightarrow$  learn to write scientifically
  - vector graphics: ipe, xfig, ... (no bitmap!)
- website will provide style files, etc.

# The Process

#### • as **author**:

write paper (5-10 pages)

• submit to conference management system

• as member of **program committee**: (

20%

10%

10%

- blind peer reviews & ranking (~2 weeks)
- back to author role:
  - submit final version  $\Rightarrow$  proceedings
  - symposium: 20 min oral presentation

# What is a data structure?

# What is a Data Structure?

A **Data Structure** specifies how to encode data from some Data Type so as to support the operators specified by a given Abstract Data Type.

# Example: Permutations

- data type: **permutation**  $\pi$  of [1,*n*]
- ADS operations:
  - $access_{\pi}(i)$ : return  $\pi[i]$
  - inverse<sub>π</sub>(j): return i such that π[i]=j
- data structure: 2 arrays A[1,n], A<sup>-1</sup>[1,n]
  - $access_{\pi}(i) = A[i]$
  - inverse $\pi(j) = A^{-1}[j]$
- might be OK, might be not (e.g. dynamic??)

# Extending Functionality

- have: DS **D** for ADT **T**
  - e.g. permutations with access/inverse
- want: DS **D'** for ADT **T'** with **T'**  $\supseteq$  **T** 
  - e.g. perms with access/inverse **plus** inversions  $\pi(i) = |\{j \le \pi[i] > \pi[i]\}|$
- use **D** as black box: **D'** is called **index** 
  - sublinear space possible:  $|\mathbf{D'}| = o(|\mathbf{D}|)$

# Implicit DS

• clever storage functionality "for free" e.g. heap:  $parent(x) = \left| \frac{x}{2} \right|$ 8 6 7 8 9 4 5 2 3 9 3 8 4



# Course Contents

- hashing
- predecessor
  data structures
- integer sorting/ searching
- distance oracles
- tree labelings

- lowest common/ level ancestors
- range minimum queries
- succinct trees
- text indexing

# Hashing

- set S of *n* objects from a LARGE universe U
- query for membership (+satellite info)
- Use space O(n), not O(|U|)



# Hashing: lookup time

- chaining/linear probing:
  O(1) expected time
- cuckoo hashing:
  O(1) worst case time
- other operations O(I) amortized & expected



### Predecessor Queries

- S: n objects from a SORTED universe U
- given  $x \in U$ , return max{ $y \le x : y \in S$ }
- fast if elements are integers: O(|g|g|U|)



# Integer Sorting

• sort *n* elements from a universe [0,2<sup>w</sup>-1]

• comparison based sorting:  $\Theta(n \lg n)$ 

- counting sort:  $O(n + 2^w)$
- with predecessor queries:  $O(n \lg w)$
- signature sort:
  - O(n) for w sufficiently large
  - $O(n \operatorname{lglg} n)$  for all w

#### Distance Oracles





# Lowest Common Ancestors



#### Level Ancestors



# Range Minimum Queries





# String **B-Trees**

- text indexing in **external** memory
- substring queries (cf suffix tree/array)
- new challanges (minimize IOs)



# Theory vs. Practice

- focus on theoretical (=mathematical) analysis of data structures
- BUT: most methods highly **practical** (perhaps with some engineering effort)
  - VL "Algorithm Engineering"
- every method better than naive approach (complex analysis ⇒ slow running time)

#### Classification of DSs



# Time vs Space

#### e.g. tree + LCA

